

# Routing Protocols for Wireless Sensor Networks

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Mohammad Hossein Yaghmaee  
Associate Professor  
Department of Computer Engineering  
Ferdowsi University of Mashhad (FUM)



# Introduction

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- Sensors are micro-electro-mechanical systems (MEMS)
- Low power devices
- Data processing capable
- Communication capabilities



# Introduction - Usage

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- Gather data locally (Temperature, Humidity, Motion Detection, etc.)
- Send them to a command center (sink)
- Applications
  - Surveillance
  - Security
  - Disaster Management
  - Environmental Studies



# Introduction - Constraints

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- Limitations
  - Energy Constrains
  - Bandwidth
- All layers must be energy aware
- Need for energy efficient and reliable network routing
- Maximize the lifetime of the network



# Differences of Routing in WSN and Traditional Networks

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- **No global addressing**
  - Classical IP-based protocols cannot be applied to sensor networks
- **Redundant data traffic**
  - Multiple sensors may generate same data within the vicinity of a phenomenon.
  - Such redundancy needs to be exploited by the routing protocols to improve energy and bandwidth utilization.



# Differences of Routing in WSN and Traditional Networks

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- **Multiple-source single-destination network**
  - Almost all applications of sensor networks require the flow of sensed data from multiple regions (sources) to a particular sink.
- **Careful resource management**
  - Sensor nodes are tightly constrained in terms of:
    - Transmission power
    - On-board energy
    - Processing capacity
    - Storage



# System architecture and design issues

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- Network Dynamics
- Node Deployment
- Energy Considerations
- Data Delivery Models
- Node capabilities
- Data aggregation/fusion



# Network Dynamics

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- Mobile or Stationary nodes
- Static Events (Temperature)
- Dynamic Events ( Target Detection)
- Dynamic events in most applications require periodic reporting and consequently generate significant traffic to be routed to the sink





# Node Deployment

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- **Deterministic :**
  - The sensors are manually placed and data is routed through pre-determined paths
- **Self-organizing :**
  - The sensor nodes are scattered randomly creating an infrastructure in an ad hoc manner.
  - Position of the sink or the cluster-head is also crucial in terms of energy efficiency and performance.



# Energy Considerations

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- Energy Considerations
  - Direct vs Multi-hop communication
    - Direct Preferred – Sensors close to sink
    - Multi-hop – unavoidable in randomly scattered networks
  - Since the transmission power of a wireless radio is proportional to distance squared
    - Multi-hop routing will consume less energy than direct communication.
  - Multi-hop routing introduces significant overhead for topology management and medium access control.
  - Direct routing would perform well enough if all the nodes were very close to the sink.



# Data Delivery Models

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- Different data delivery models depending on the application of the sensor network:
  - Continuous
    - Each sensor sends data periodically
  - Event-driven:
    - The transmission of data is triggered when an event occurs.
  - Query-driven:
    - The transmission of data is triggered when a query is generated by the sink
  - Hybrid
- The routing protocol is highly influenced by the data delivery model



# Node Capabilities

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- In a sensor network, different functionalities can be associated with the sensor nodes.
  - Homogenous:
    - All nodes have equal capacity in terms of computation, communication and power
  - Heterogeneous
    - Nodes dedicated to a particular task (relaying, sensing, aggregation)



# Data Aggregation/Fusion

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- Similar packets from multiple nodes can be aggregated
  - The number of transmissions would be reduced.
- Data aggregation
  - Combination of data using functions such as suppression (eliminating duplicates), min, max and average
  - Aggregation – Combination of data by eliminating redundancy
  - Data Fusion is Aggregation through signal processing techniques
  - Aggregation achieves energy savings



# Classification of Routing Protocols

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- Data Centric:
  - Data-centric protocols are query-based
- Hierarchical:
  - Aim at clustering the nodes so that cluster heads can do some aggregation and reduction of data in order to save energy
- Location-based:
  - Utilize the position information to relay the data to the desired regions rather than the whole network.
- Network Flow & QoS Aware:
  - Are based on general network-flow modeling and protocols that strive for meeting some QoS requirements along with the routing function



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# Data-centric Protocols

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- In many applications of sensor networks, it is not feasible to assign global identifiers to each node
- Data-centric protocols are query-based.
- Sink sends queries to certain regions and waits data from sensors located in that region
- Attribute-based naming is necessary to specify properties of data





# Data-centric Routing

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- Sensor networks can be considered as a virtual database
- Implement query-processing operators in the sensor network
- Queries are flooded through the network (or sent to a representative set of nodes)
- In response, nodes generate tuples and send matching tuples towards the origin of the query
- Intermediate nodes may merge responses or aggregate



# Data-centric Protocols

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- Flooding
- Gossiping
- Sensor Protocols for Information via Negotiation (SPIN)
- Directed Diffusion
- Energy-aware Routing
- Rumor Routing
- Gradient-Based Routing (GBR)
- Constrained Anisotropic Diffusion Routing (CADR)
- COUGAR
- Active QUery forwarding In sensoR nEtworks (ACQUIRE)



# Data-centric Protocols

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- **Flooding**

- Sensor broadcasts every packet it receives
- Relay of packet till the destination or maximum number of hops
- No topology maintenance or routing

- **Gossiping**

- Enhanced version of flooding
- Sends received packet to a randomly selected neighbor

# Classic Flooding Problems

## ■ **Implosion Problem:**

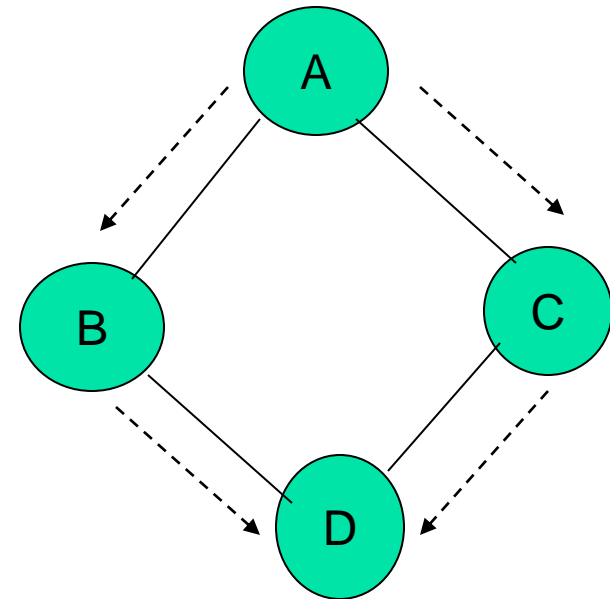
- A starts by flooding its data to all of its neighbors.
- Two copies of the data eventually end at node D.
- The system wastes energy and bandwidth.

## ■ **Overlap Problem:**

- Two sensors cover an overlapping graphic region.
- Node receives two copies of the Data.

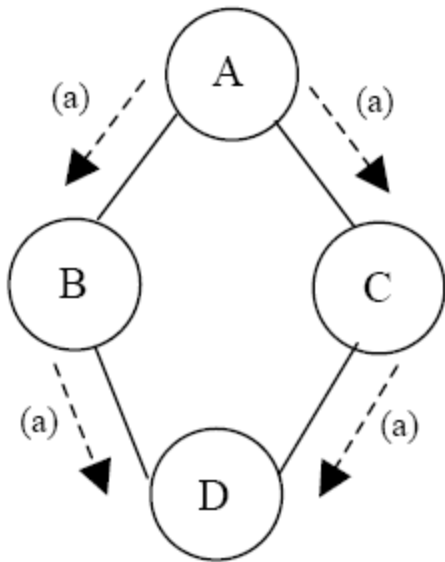
## ■ **Resource Blinding:**

- Resources do not modify their activities based on the amount of energy they have.

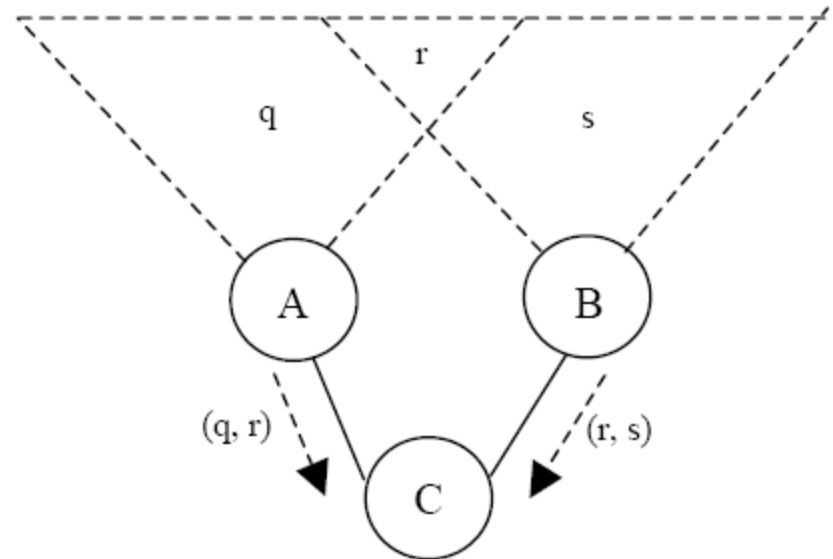


# Data-centric Protocols – Flooding, Gossiping Problems

- Problems of Implosion, Overlap, Resource Blindness



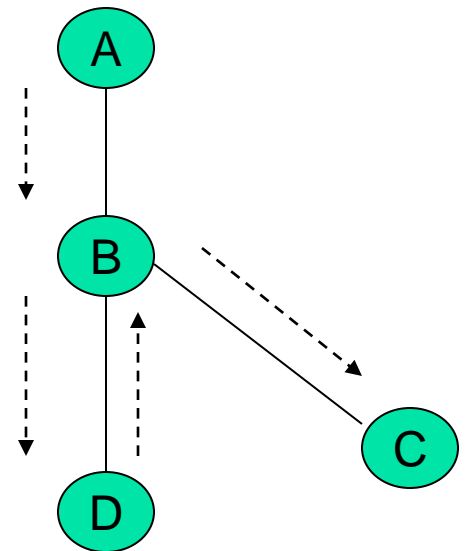
**Implosion Problem**



**Overlap Problem**

# Gossiping

- An alternative to the classic flooding
- Uses randomization to conserve energy.
- Each node only forwards data on to one neighbor
  - Is selected randomly.
- After node D receives the data, it must forward the data back to sender (B)
  - Otherwise the data would never reach node C



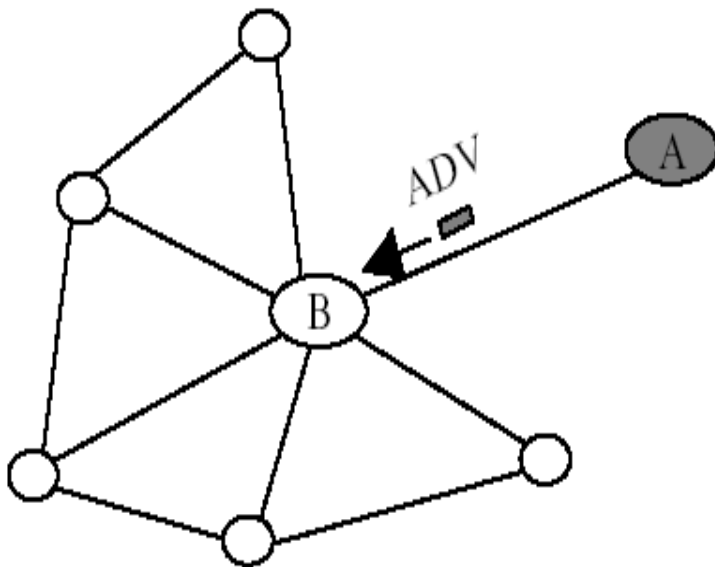


# SPIN: Sensor Protocols for Information Negotiation

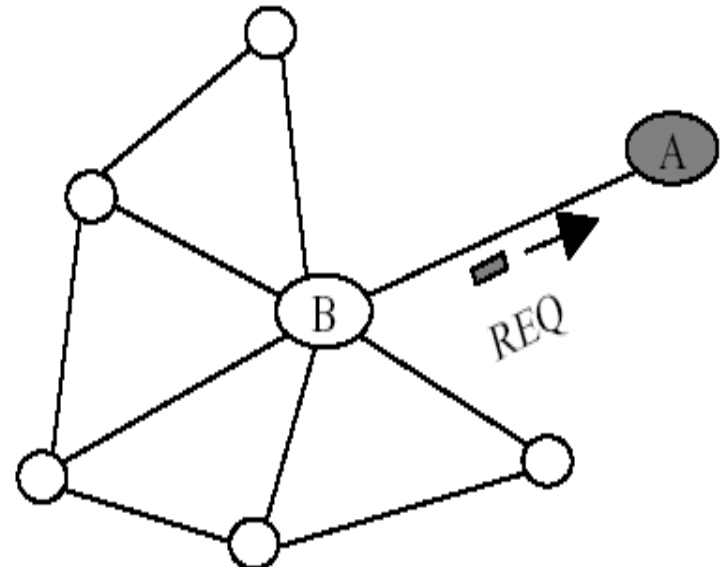
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- One of the most dominant form of routing in the wireless sensor networks.
- Name data, using meta-data
  - Meta Data for each sensor data
  - Same sensor data -> same meta-data
  - Different sensor data -> different meta-data
- Size of meta-data << Size of actual data
- There is no standard meta-data format and it is assumed to be application specific
- Uses three types of messages:
  - ADV – advertise data
  - REQ – request for data
  - DATA – data message, contains actual sensor data

# SPIN Protocol Example



(a)

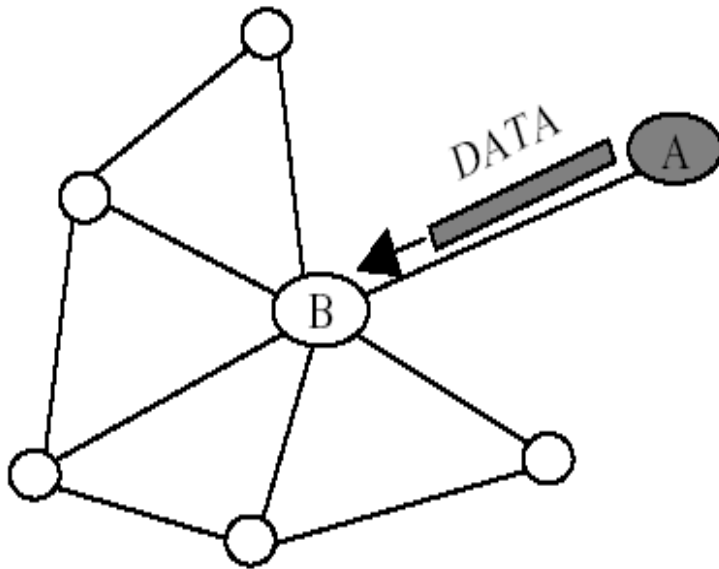


(b)

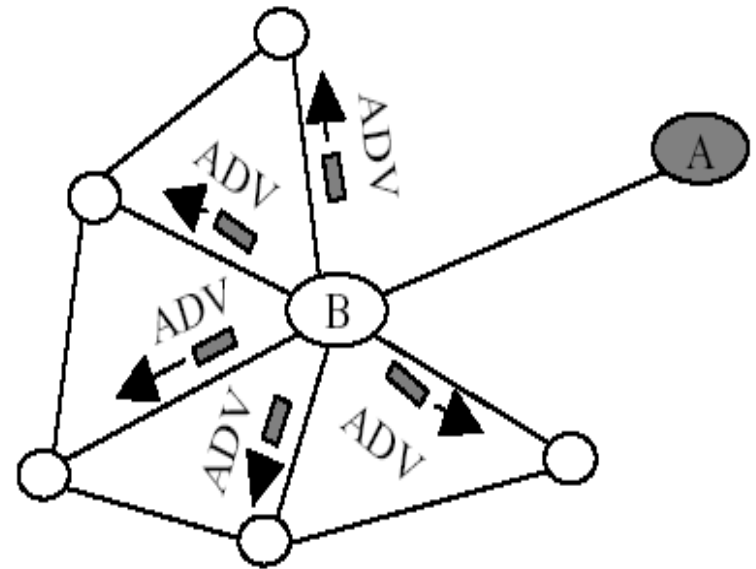
- A sends an ADV message to B
- B sends a REQ listing all of the data it would like to acquire.



# SPIN Protocol Example



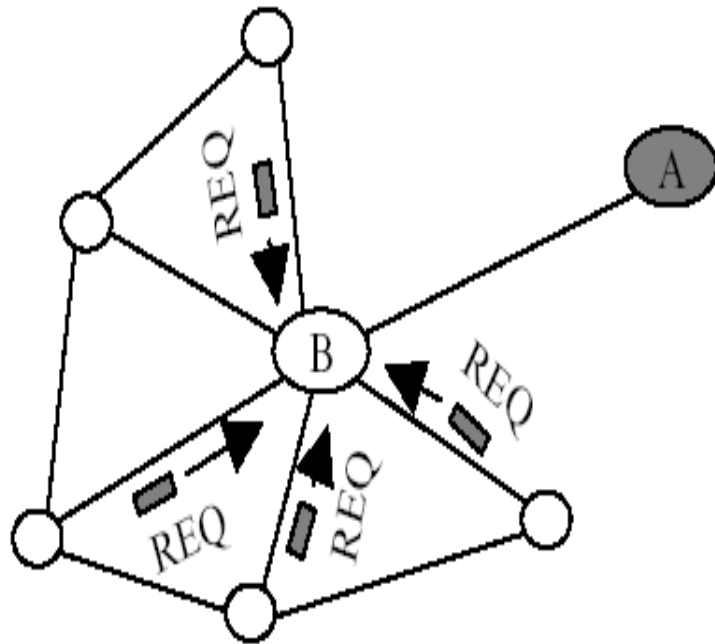
(c)



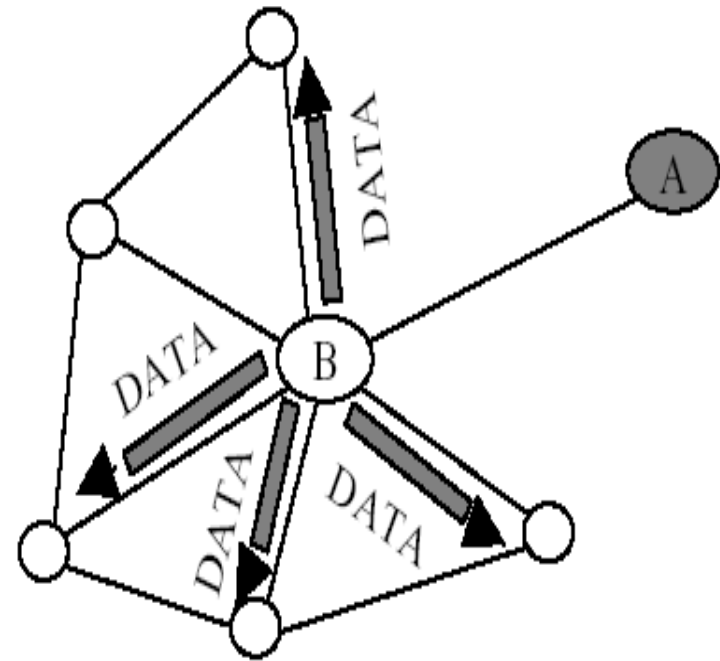
(d)

- If node B had its own data, it could aggregate this with the data of node A and advertise.

# SPIN Protocol Example



(e)



(f)

Nodes need not respond to every message



# Data-centric Protocols - SPIN

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- Topological changes are localized - Each node needs to know only its neighbors
- SPIN halves the redundant data in comparison to flooding
- Cannot guarantee data delivery
- SPIN **NOT** good for applications that need reliable data delivery



# SPIN1 and SPIN2

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**SPIN1** : Three way handshaking protocol.

- ADV, REQ, DATA.
- Each sensor node has resource manager
  - Keeps track of resource consumption
  - Applications probe the manager before any activity
  - Cut down activity to save energy

**SPIN2** : energy constraint

- Adds energy-conservative heuristic to the SPIN1 protocol.
- Node initiates three stage protocol, only if it has enough energy to complete it.
- If below energy threshold, node can still receive messages, cannot send/recv DATA messages



# Direct Diffusion: Motivation

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- Properties of Sensor Networks
  - Data centric
  - No central authority
  - Resource constrained
  - Nodes are tied to physical locations
  - Nodes may not know the topology
  - Nodes are generally stationary
- How can we get data from the sensors?



# Direct Diffusion

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- Uses a naming scheme for the data to save energy
- Attribute-value pairs for data and queries on-demand (Interests)
- Interests are broadcasted by the sink (query) to its neighbors (caching), which can do in-network aggregation
- Gradients = reply links to an interest (path establishment)



# Direct Diffusion

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- Direct Diffusion suggests the use of attribute-value pairs for the data and queries the sensors in an on demand basis by using those pairs.
- In order to create a query, an interest is defined using a list of attribute-value pairs such as:
  - name of objects,
  - interval,
  - duration,
  - geographical area,
  - etc.

# Directed Diffusion: Main Features



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- Data centric
  - Individual nodes are unimportant
- Request driven
  - Sinks place requests as interests
  - Sources satisfying the interest can be found
  - Intermediate nodes route data toward sinks
- Localized repair and reinforcement
- Multi-path delivery for multiple sources, sinks, and queries





# Direct Diffusion

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- The interest is broadcast by a sink through its neighbors.
- Each node receiving the interest can do caching for later use.
- The nodes also have the ability to do in-network data aggregation
- The interests in the caches are then used to compare the received data with the values in the interests.



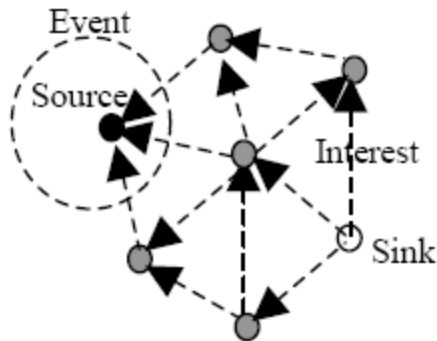
# Direct Diffusion

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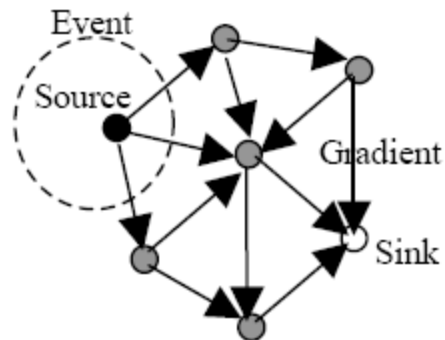
- The interest entry also contains several gradient fields.
- A gradient is a reply link to a neighbor from which the interest was received.
- It is characterized by the data rate, duration and expiration time derived from the received interests fields.
- By utilizing interest and gradients, paths are established between sink and sources.
- Several paths can be established so that one of them is selected by reinforcement.
- The sink resends the original interest message through the selected path with a smaller interval hence reinforces the source node on that path to send data more frequently

# Data-centric Protocols – Direct Diffusion

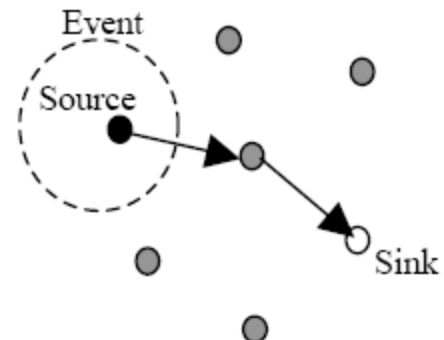
- Energy saving and delay done with caching
- No need for global addressing (neighbor-to-neighbor mechanism)
- **Cannot** be used for continuous data delivery or event-driven applications



(a) Interest propagation



(b) Initial gradients setup



(c) Data delivery along reinforced



# Directed Diffusion: Motivating Example

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- Sensor nodes are monitoring animals
- Users are interested in receiving data for all 4-legged creatures seen in a rectangle
- Users specify the data rate



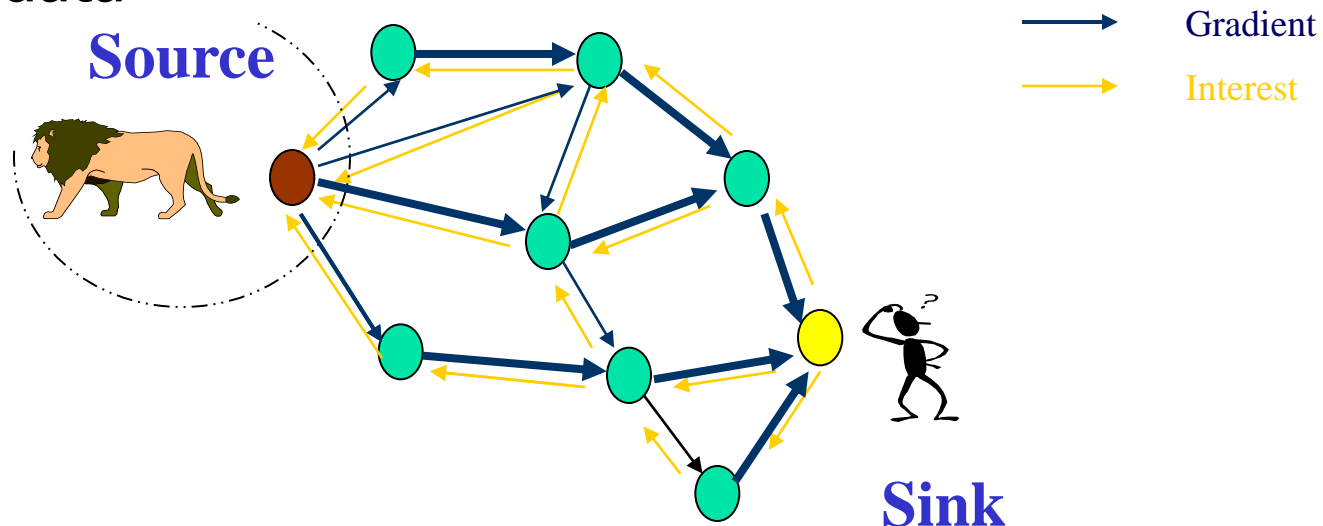
# Directed Diffusion: Interest and Event Naming

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- Query/interest:
  1. Type=four-legged animal
  2. Interval=20ms (event data rate)
  3. Duration=10 seconds (time to cache)
  4. Rect=[-100, 100, 200, 400]
- Reply:
  1. Type=four-legged animal
  2. Instance = elephant
  3. Location = [125, 220]
  4. Intensity = 0.6
  5. Confidence = 0.85
  6. Timestamp = 01:20:40
- Attribute-Value pairs, no advanced naming scheme

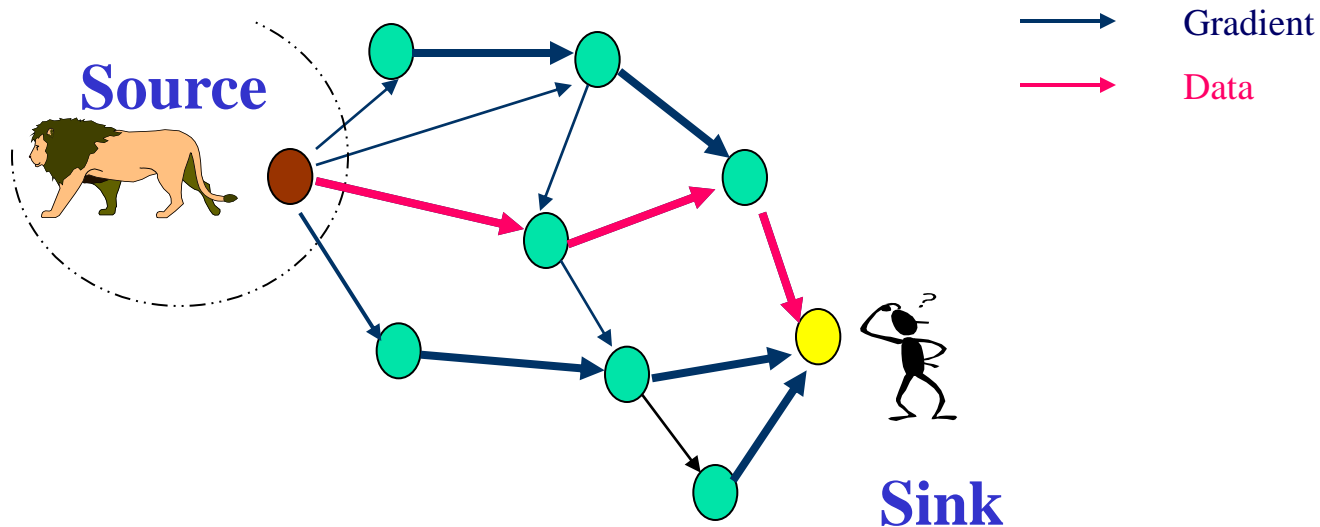
# Directed Diffusion: Interest Propagation

- Flood interest
- Constrained or Directional flooding based on location is possible
- Directional propagation based on previously cached data



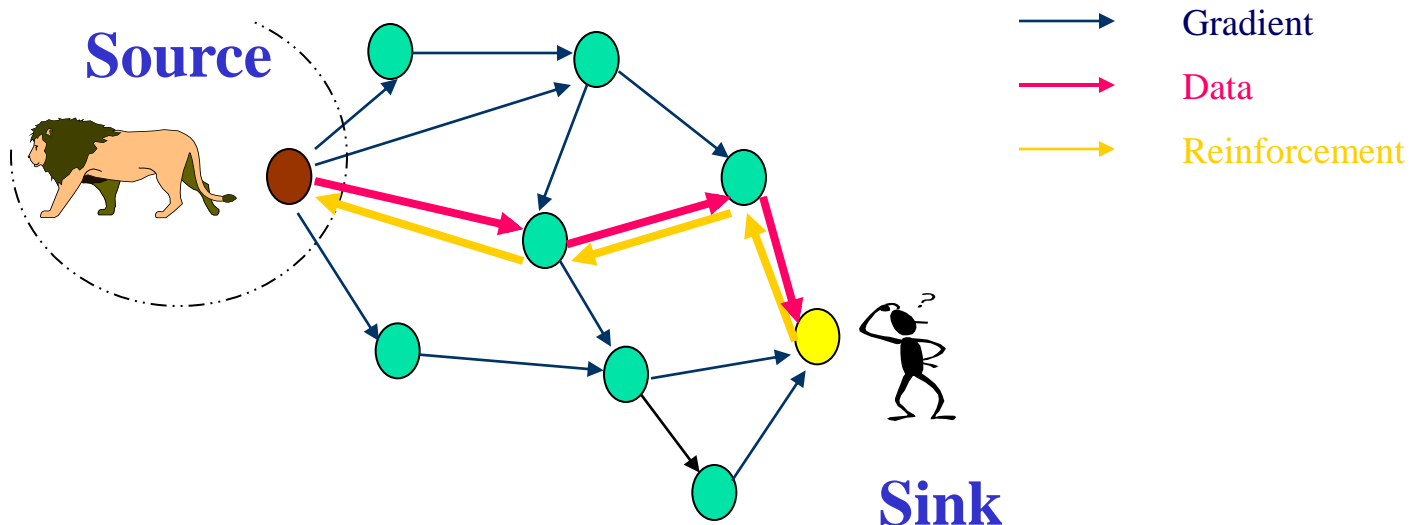
# Directed Diffusion: Data Propagation

- Multipath routing
  - Consider each gradient's link quality



# Directed Diffusion: Reinforcement

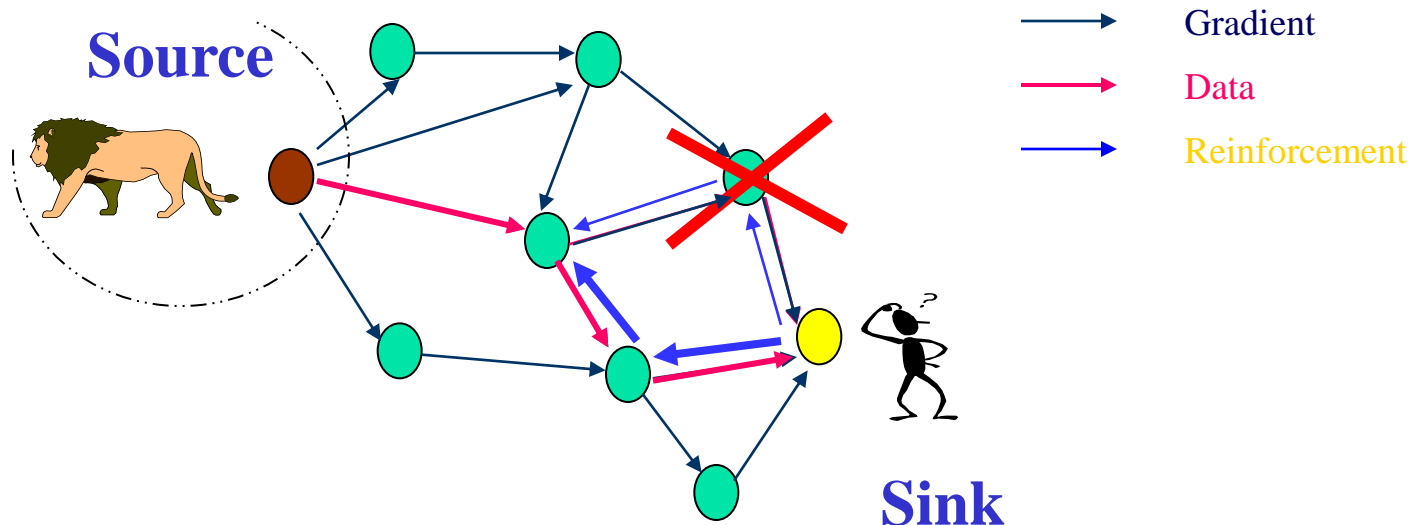
- Reinforce one of the neighbor after receiving initial data.
  - Neighbor who consistently performs better than others
  - Neighbor from whom most events received





# Directed Diffusion: Negative Reinforcement

- Explicitly degrade the path by re-sending *interest* with lower data rate.
- Time out: Without periodic reinforcement, a gradient will be torn down





# Directed Diffusion Conclusion

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- Different from SPIN in terms of on-demand data querying mechanism
  - Sink floods interests only if necessary
    - A lot of energy savings
  - In SPIN, sensors advertise the availability of data
- Characteristics
  - Data centric:
    - All communications are neighbor to neighbor with no need for a node addressing mechanism
  - Each node can do aggregation & caching
  - On-demand, query-driven:
    - Inappropriate for applications requiring continuous data delivery, e.g., environmental monitoring
  - Attribute-based naming scheme is application dependent
    - For each application it should be defined a priori
    - Extra processing overhead at sensor nodes



# Energy-aware Routing

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- Chose paths base on a probability function, which depends on the energy consumption of each path.
- Network survivability is the main metric
- Using the minimum energy path all the time will deplete the energy of nodes on that path.
- Instead, one of the multiple paths is used with a certain probability so that the whole network lifetime increases.
- The protocol assumes that each node is addressable through a class-based addressing which includes the location and types of the nodes.



# Energy-aware Routing

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- Occasional use of a set of sub-optimal paths
- Multiple paths used with certain probability
- Increase of the total lifetime of the network
- Hinders the ability for recovering from node failure
- Requires address mechanism → Complicate setup



# Energy-aware routing

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- There are 3 phases in the protocol.
  - Setup phase:
    - Localized flooding occurs to find the routes and create the routing tables.
  - Data communication phase:
    - Each node forwards the packet by randomly choosing a node from its forwarding table using the probabilities.
  - Route maintenance phase:
    - Localized flooding is performed infrequently to keep all the paths alive.



# Setup Phase (Cost Function)

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- If the request is sent from node  $N_i$  to node  $N_j$ ,  $N_j$  calculates the cost of the path as follows:

$$C_{N_j, N_i} = \text{Cost}(N_i) + \text{Metric}(N_j, N_i)$$

- Paths that have a very high cost are discarded
- The node assigns a probability to each of its neighbors in routing (forwarding) table (FT) corresponding to the formed paths
- The probability is inversely proportional to the cost

$$P_{N_j, N_i} = \frac{1/C_{N_j, N_i}}{\sum_{k \in \text{FT}_j} 1/C_{N_j, N_k}}$$



# Setup Phase (Cost Function)

---

- $N_j$  then calculates the average cost for reaching the destination using the neighbors in the forwarding table ( $FT_j$ ) using the formula:

$$\text{Cost}(N_j) = \sum_{i \in FT_j} P_{N_j, N_i} C_{N_j, N_i}$$

- This average cost for  $N_j$  is set in the cost field of the request and forwarded.



# Comparison with Directed Diffusion

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- Like Directed Diffusion, potential paths from data sources to the sink are discovered.
- In Directed Diffusion, data is sent through multiple paths, one of them being reinforced to send at higher rates.
- Energy aware routing protocol selects a single path randomly from the multiple alternatives in order to save energy.
- It provides an overall improvement of 21.5% energy saving and a 44% increase in network lifetime.
- It hinders the ability of recovering from a node or path failure as opposed to Directed Diffusion.
- It requires gathering the location information and setting up the addressing mechanism for the nodes.
- More complicate route setup compared to the Directed Diffusion.





# Rumor Routing

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- Variation of Directed Diffusion
- Flood the events instead of the queries
- Creation of an event → generation of a long live packet travel through the network (agent)
- Nodes save the event in a local table
- When a node receives query → checks its table and returns source – destination path



# Rumor Routing

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- Variation of directed diffusion
  - Don't flood interests (or queries)
  - Flood events when the number of events is small but the number of queries large
  - Route the query to the nodes that have observed a particular event
  - Long-lived packets, called agents, flood events through the network
  - When a node detects an event, it adds the event to its events table, and generates an agent
  - Agents travel the network to propagate info about local events
    - An agent is associated with TTL (Time-To-Live)



# Rumor Routing

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- Advantages
  - Can handle node failure
  - Significant energy savings
- Disadvantages
  - Works well **only** with small number of events
  - Overhead through adjusting parameters, like the time to live of the agent



# Gradient-Based Routing (GBR)

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- Slightly changed version of Directed Diffusion
- Keep the number of hops when the interest is diffused through the network.
- Hence, each node can discover the minimum number of hops to the sink, which is called height of the node.
- The difference between a nodes height and that of its neighbor is considered the gradient on that link.
- A packet is forwarded on a link with the largest gradient.



# GBR

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- It uses traffic spreading and data aggregation to balance uniformly the network traffic
- The data spreading schemes strives to achieve an even distribution of the traffic throughout the whole network.
  - helps in balancing the load on sensor nodes and increases the network lifetime.
- Outperforms Directed Diffusion in terms of total communication energy



# Traffic Spreading Methods in GBR

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- Stochastic scheme:
  - When there are two or more next hops with the same gradient, the node chooses one of them at random.
- Energy-based scheme:
  - When a nodes energy drops below a certain threshold, it increases its height so that other sensors are discouraged from sending data to that node.
- Stream-based scheme:
  - The idea is to divert new streams away from nodes that are currently part of the path of other streams.



# Constrained Anisotropic Diffusion Routing (CADR)

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- A general form of Directed Diffusion.
- The idea is to query sensors and route data in a network in order to maximize the information gain, while minimizing the latency and bandwidth.
- This is achieved by activating only the sensors that are close to a particular event and dynamically adjusting data routes.
- The major difference from Directed Diffusion is the consideration
- of information gain in addition to the communication cost.
- Each node evaluates an information/cost objective and routes data based on the local information/cost gradient and end-user requirements.



# Constrained Anisotropic Diffusion Routing (CADR)

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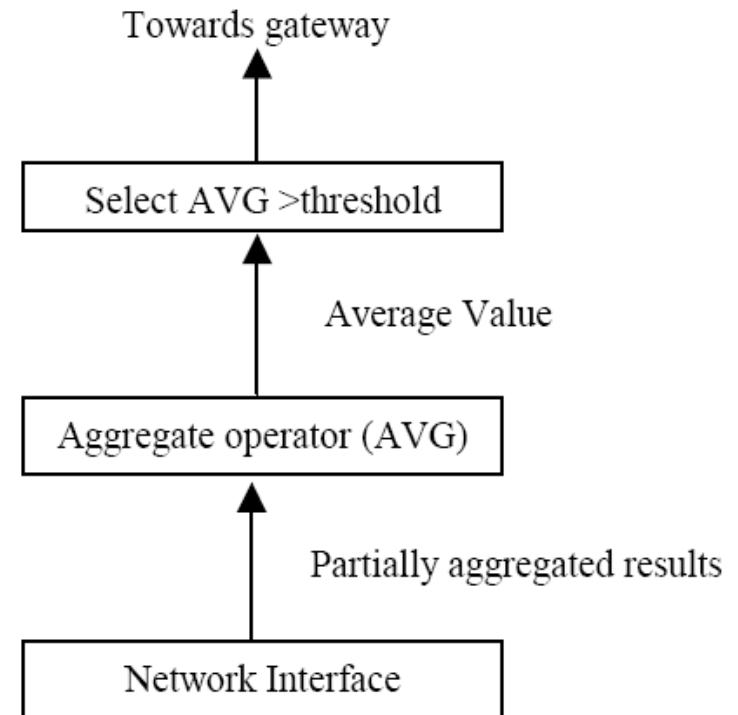
- General form of Directed Diffusion
- Query Sensors
- Route data in the network
- Activates sensors close to the event and dynamically adjusts routes
- Routing based on a local information/cost gradient
- More energy efficient than Directed Diffusion



# Data-centric Protocols

## ■ COUGAR

- Views the network as a huge distributed database
- Declarative queries to abstract query processing from network layer functions
- Introduces a new query layer
- Leader node performs data aggregation and transmits to the sink



Query plan at a leader node



# Data-centric Protocols - COUGAR

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- Disadvantages

- Additional query layer brings overhead in terms of energy consumption and storage
- In network data computation requires synchronization (i.e. wait for all data before sending data)
- Dynamically maintenance of leader nodes to prevent failure



# Active QUery forwarding In sensor nEtworks (ACQUIRE)

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- Views network as a distributed database
- Node receiving a query from the sink tries to respond partially and then forwards packet to a neighbor
- Use of pre-cached information
- After the query is answered, result is returned to the sink by using the reverse path or the shortest path
- If cache information is not up to date → node gathers information from neighbors within look ahead of  $d$  hops



# ACQUIRE

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- Motivation: Deal with one shot complex queries
- Efficient routing by adjusting parameter  $d$
- If  $d$  equals network size  $\rightarrow$  behaves similar to flooding
- If  $d$  too small the query has to travel more hops



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# Hierarchical Routing Protocols

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- Scalability is one of the major design attributes of sensor networks.
- A single-tier network can cause the gateway to overload with the increase in sensors density
  - Such overload might cause latency in communication and inadequate tracking of events.
- The single-gateway architecture is not scalable for a larger set of sensors covering a wider area of interest.



# Hierarchical Protocols

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- Maintain energy consumption of sensor nodes
  - By multi-hop communication within a particular cluster
  - By data aggregation and fusion → decrease the number of the total transmitted packets



# Hierarchical Protocols

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- LEACH : Low-Energy Adaptive Clustering Hierarchy
- PEGASIS: Power-Efficient GATHERing in Sensor Information Systems
  - Hierarchical PEGASIS
- TEEN: Threshold sensitive Energy Efficient sensor Network protocol
  - Adaptive Threshold TEEN (APTEEN)
- Energy-aware routing for cluster-based sensor networks
- Self-organizing protocol





# LEACH : Low-Energy Adaptive Clustering Hierarchy

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- One of the first hierarchical routing protocols
- Forms clusters of the sensor nodes based on received signal strength
- Local cluster heads route the information of the cluster to the sink
- Cluster heads change randomly over time → balance energy dissipation
- Data processing & aggregation done by cluster head



# Cluster Head (CH)

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- Each node randomly decides to become a cluster heads (CH)
- CH chooses the code to be used in its cluster
  - CDMA between clusters
- CH broadcasts Adv;
  - Each node decides to which cluster it belongs based on the received signal strength of Adv
- CH creates a transmission schedule for TDMA in the cluster
  - Nodes can sleep when its not their turn to transmit
- CH compresses data received from the nodes in the cluster and sends the aggregated data to BS
- CH is rotated randomly



# Cluster Head Choosing

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- All the data processing such as data fusion and aggregation are local to the cluster.
- CHs change randomly over time in order to balance their energy dissipation of nodes.
- This decision is made by the node choosing a random number between 0 and 1.
- The node becomes a cluster head for the current round if the number is less than a threshold.



# LEACH Threshold Function

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$$T(n) = \begin{cases} \frac{p}{1 - p * (r \bmod 1/p)} & \text{if } n \in G, \\ 0 & \text{otherwise,} \end{cases}$$

- $p$  is the desired percentage of cluster heads (e.g. 0.05),
- $r$  is the current round,
- $G$  is the set of nodes that have not been cluster heads in the last  $1/p$  rounds.



# LEACH Conclusion

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- **Advantages**

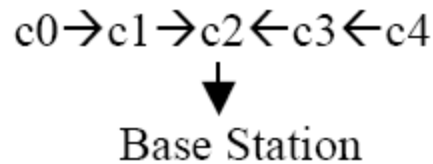
- Completely distributed
- No global knowledge of the network
- Increases the lifetime of the network

- **Disadvantages**

- Uses single-hop routing within cluster → **not** applicable to networks in large regions
- Dynamic clustering brings extra overhead (advertisements, etc)

# PEGASIS: Power-Efficient GATHERing in Sensor Information Systems

- Improvement of LEACH
- Forms chains from sensors rather than clusters



- Data aggregation in the chain → one node sends the data to the base station
- Outperforms LEACH
- Excessive delay for distant nodes in the chain



# PEGASIS Characteristics

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- Use multi-hop routing by forming chains.
- Selecting only one node to transmit to the base station instead of using multiple nodes.
- PEGASIS has been shown to outperform LEACH by about 100–300% for different network sizes and topologies.
- PEGASIS introduces excessive delay for distant node on the chain.
- The single leader can become a bottleneck.



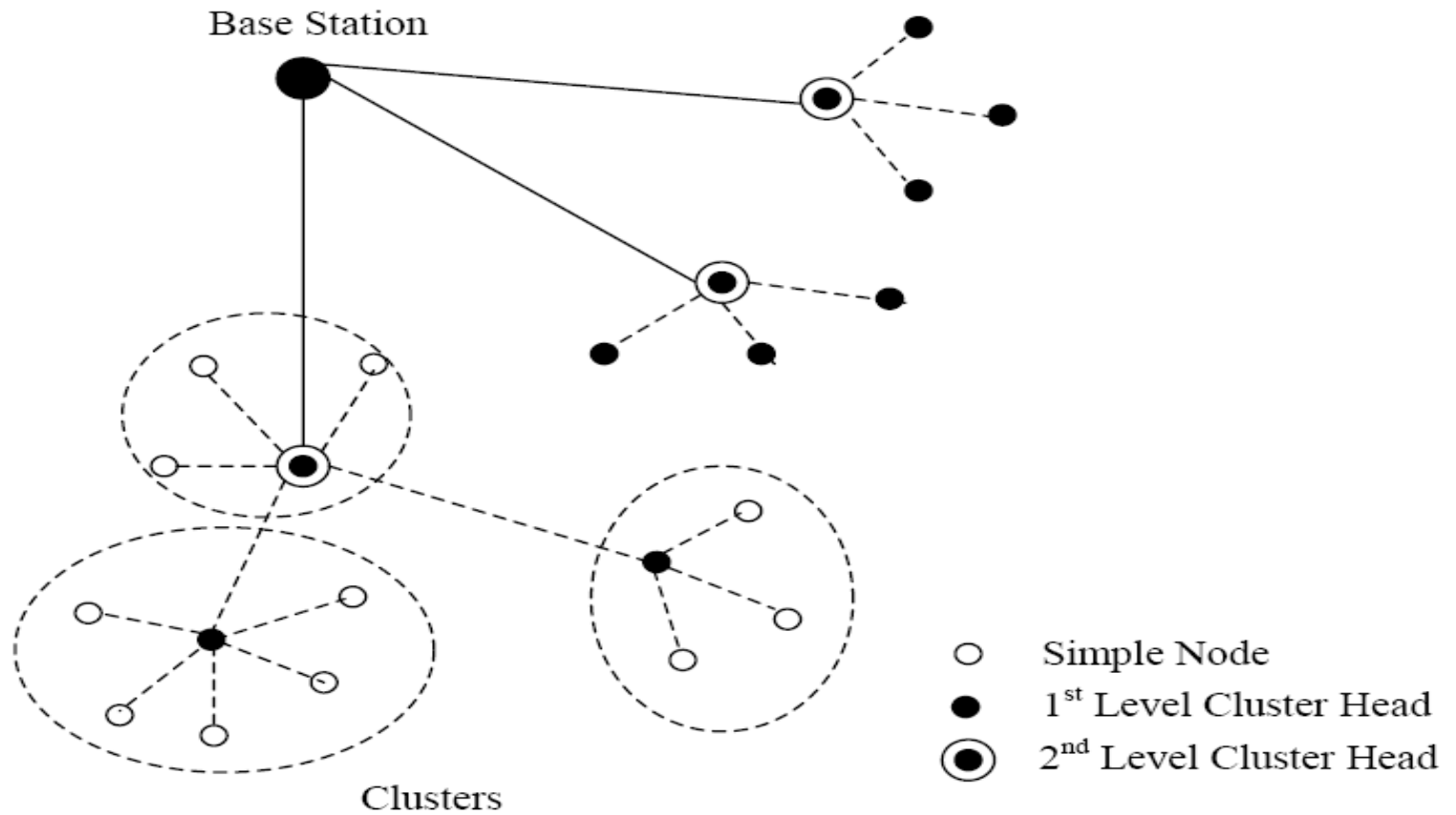
# TEEN: Threshold sensitive Energy Efficient Sensor Network Protocol

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- Good for time-critical applications
- Hierarchical along with a data-centric approach
- Hierarchical grouping:
  - Close nodes form clusters and this process goes on the second level until sink is reached
- Not good for applications that need periodic reports



# TEEN Example





# TEEN Thresholds

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- Cluster headers broadcast:
  - Hard Threshold
    - the minimum possible value of an attribute to trigger a sensor node to transmit to the cluster head
    - reducing the number of transmissions significantly
  - Soft Threshold
    - Once a node senses a value at or beyond the hard threshold, it transmits data only when the value of that attribute changes by an amount equal to or greater than the soft threshold



# TEEN Conclusion

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- Advantages

- Outperform LEACH in terms of energy dissipation and total lifetime of the network

- Disadvantages

- Overhead and complexity of:
  - Forming multiple level clusters
  - Implementing threshold-based functions
  - Dealing with attribute-based naming of queries

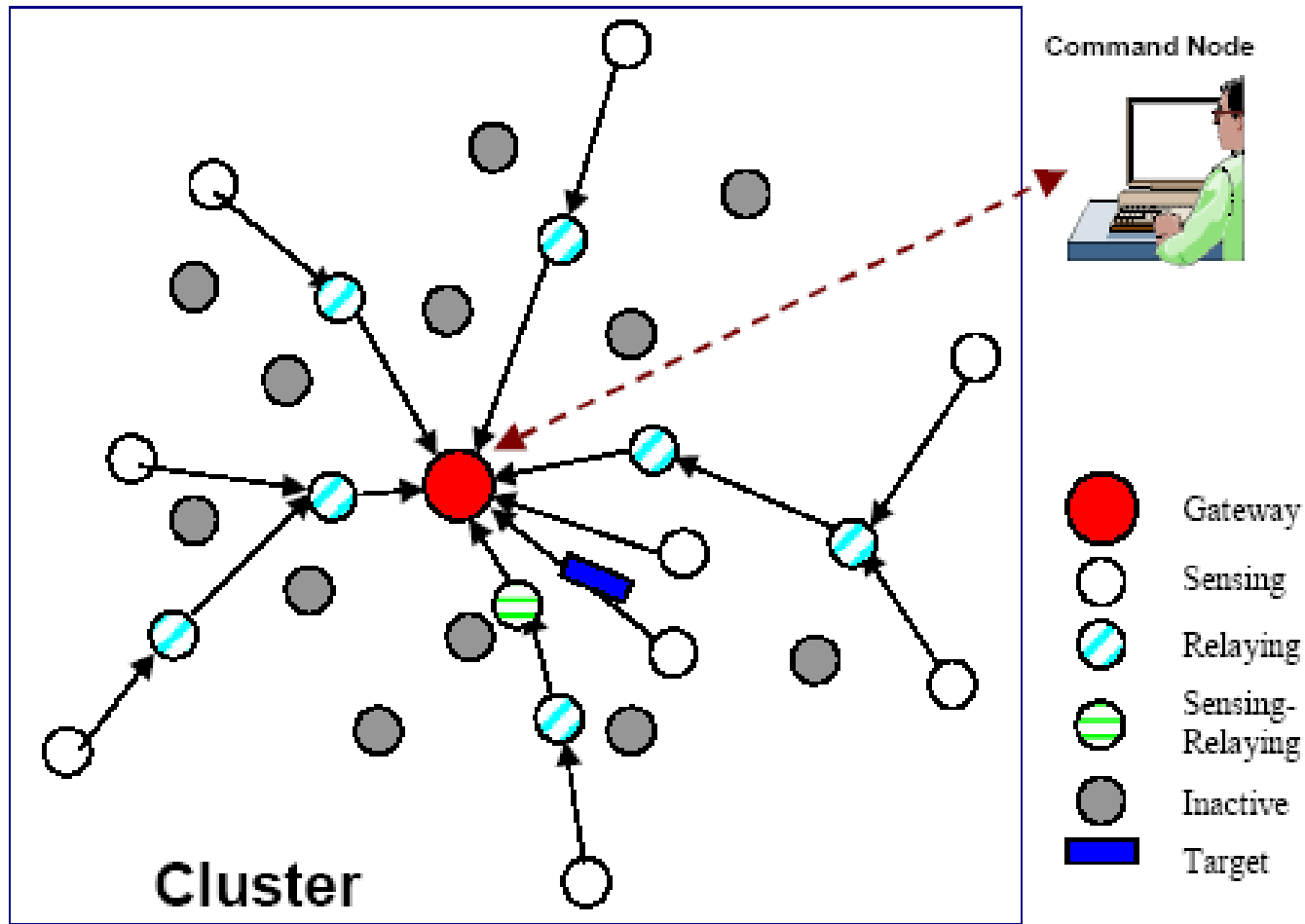


# Energy-aware Routing For Cluster-based Sensor Networks

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- Assumptions:
  - Sensors are grouped into clusters prior to network operation
  - Cluster Heads (Gateways) less energy constrained
  - Cluster Heads know the location of the sensors → Known Multi-Hop routing to collect data
  - Communication node (sink) communicates only with gateways

# Architecture





# Stages of a Sensor

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- Stages of a Sensor inside a cluster
  - Sensing only:
    - the node probes the environment and generates data at a constant rate.
  - Relaying only :
    - the node does not sense the target but its communications circuitry is on to relay the data from other active nodes.
  - Sensing-Relaying:
    - node is both sensing and relaying messages from other nodes
  - Inactive:
    - the node can turn off its sensing and communication circuitry.



# Self-organizing Protocol

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- The architecture supports heterogeneous sensors that can be mobile or stationary
- Sensors probe the environment and forward the data to designated routers.
- Router nodes are stationary and form the backbone for communication.
- Collected data are forwarded through the routers to more powerful sink nodes.



# Self-organizing Protocol

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- The architecture requires addressing
  - Sensor identified by the router is connected to
- Sensing nodes are identifiable through the address of the router node it is connected to.
- The routing architecture is hierarchical where groups of nodes are formed and merge when needed.
- Utilizes router nodes to keep all sensors connected by forming a dominating set





# Different Phases

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- Discovery phase:
  - The nodes in the neighborhood of each sensor are discovered.
- Organization phase:
  - Groups are formed and merged by forming a hierarchy.
  - Each node is allocated an address based on its position in the hierarchy.
  - Routing tables are created for each node.
  - Broadcast trees that span all the nodes are constructed.
- Maintenance phase:
  - Updating of routing tables and energy levels of nodes is made.
  - Each node informs the neighbors about its routing table and energy level.
- Self-reorganization phase:
  - In case of partition or node failures, group reorganizations are performed.



# Classification of Routing Protocols

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- Data Centric:
  - Data-centric protocols are query-based
- Hierarchical:
  - Aim at clustering the nodes so that cluster heads can do some aggregation and reduction of data in order to save energy
- Location-based:
  - Utilize the position information to relay the data to the desired regions rather than the whole network.
- Network Flow & QoS Aware:
  - Are based on general network-flow modeling and protocols that strive for meeting some QoS requirements along with the routing function



# Location-based Protocols

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- Most of the routing protocols for sensor networks require location information for sensor nodes.
- There is no addressing scheme for sensor networks like IP-addresses
- location information can be utilized in routing data in an energy efficient way.
- Protocols designed for Ad hoc networks with mobility in mind
  - Applicable to Sensor Networks as well
  - Only energy-aware protocols are considered.



# Location-based Protocols

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- **MECN & SMECN**
  - Minimum Energy Communication Network
  
- **GAF**
  - Geographic Adaptive Fidelity
  
- **GEAR**
  - Geographic and Energy Aware Routing



# MECN & SMECN

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- Utilizes low power GPS
- Best applicable to non-mobile sensor networks
- Identifies a **relay region** for every node
  - The relay region consists of nodes in a surrounding area where transmitting through those nodes is more energy efficient than direct transmission.
- The main idea of MECN is to find a sub-network, which will have less number of nodes and require less power for transmission between any two particular nodes
- Self-reconfiguring
- Dynamically adaptive

# GAF: Geographic Adaptive Fidelity

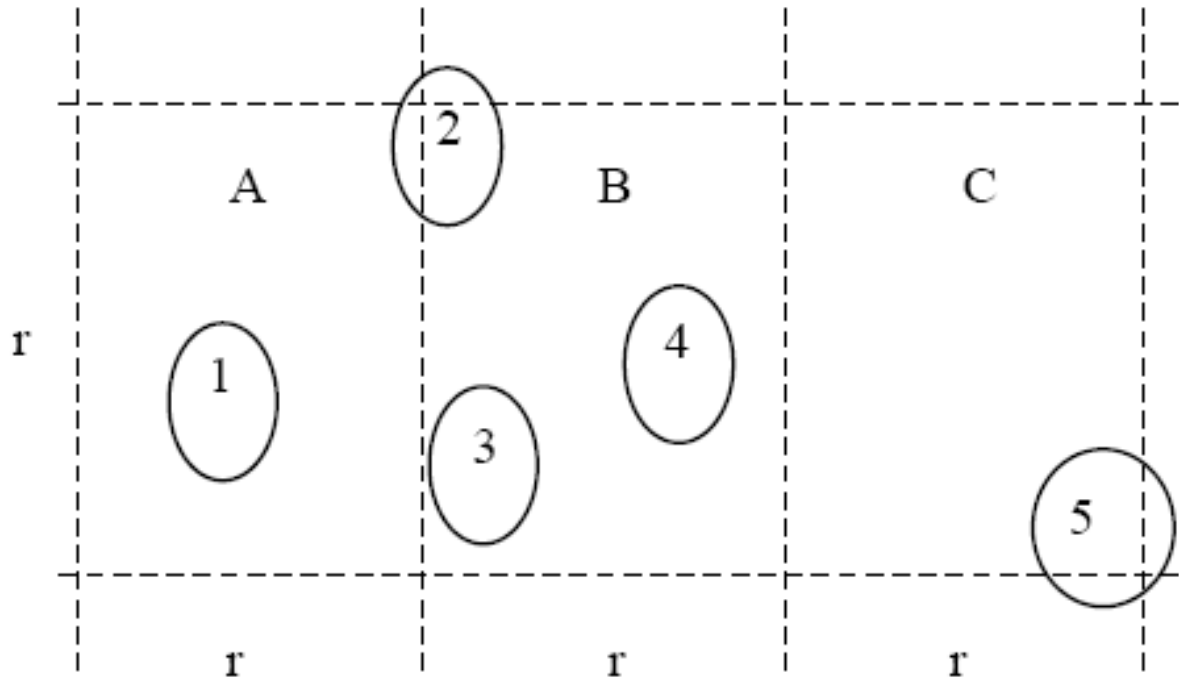


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- GAF is an energy-aware location-based routing algorithm.
- GAF conserves energy by turning off unnecessary nodes in the network without affecting the level of routing fidelity.
- It forms a virtual grid for the covered area.
- Each node uses its GPS-indicated location to associate itself with a point in the virtual grid.
- Nodes associated with the same point on the grid are considered equivalent in terms of the cost of packet routing.

# GAF Example

- Node 1 can reach any of 2, 3 and 4 and nodes 2, 3, and 4 can reach 5.
- Therefore nodes 2, 3 and 4 are equivalent and two of them can sleep.





# GAF States

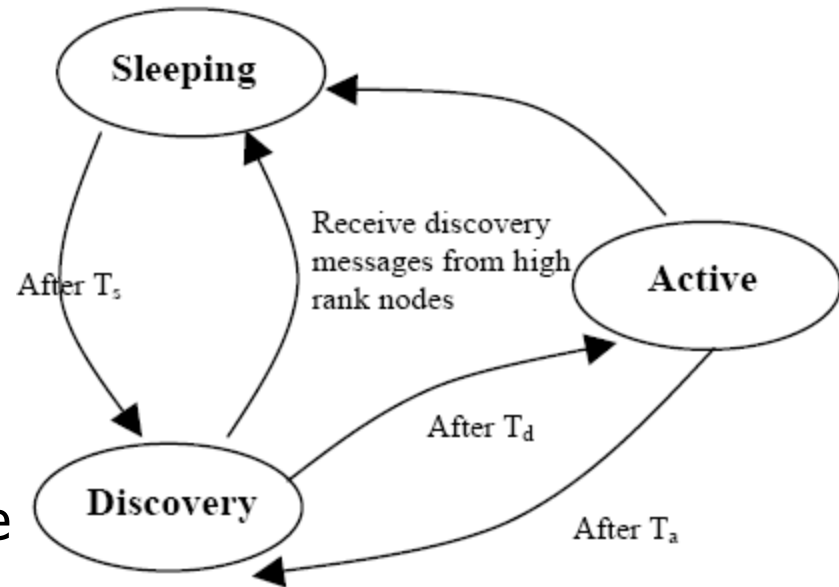
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- Three States
  - Discovery
  - Active
  - Sleep
- Discovery state is used for determining the neighbors in the grid.
- Nodes change states from sleeping to active in turn so that the load is balanced.
- Active reflecting participation in routing and sleep when the radio is turned off.
- As good as a normal Ad hoc in terms of latency and packet loss (saving energy)



# GAF State Diagram

- Each node in the grid estimates its leaving time of grid and sends this to its neighbors.
- The sleeping neighbors adjust their sleeping time accordingly in order to keep the routing fidelity.
- Before the leaving time of the active node expires, sleeping nodes wake up and one of them becomes active.
- GAF strives to keep the network connected by keeping a representative node always in active mode for each region on its virtual grid.





# Classification of Routing Protocols

---

- Data Centric:
  - Data-centric protocols are query-based
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  - Aim at clustering the nodes so that cluster heads can do some aggregation and reduction of data in order to save energy
- Location-based:
  - Utilize the position information to relay the data to the desired regions rather than the whole network.
- Network Flow & QoS Aware:
  - Are based on general network-flow modeling and protocols that strive for meeting some QoS requirements along with the routing function

# Network Flow & QoS-aware Protocols



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- Network Flow:
  - Maximize traffic flow between two nodes, respecting the capacities of the links
- QoS-aware protocols:
  - Consider end-to-end delay requirements while setting up paths

# Network Flow & QoS-aware Protocols



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- Maximum Lifetime Energy Routing
- Maximum Lifetime Data Gathering
- Minimum Cost Forwarding
- Sequential Assignment Routing
- Energy Aware QoS Routing Protocol
- SPEED

# Maximum Lifetime Energy Routing



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- Maximizes network lifetime by defining link cost as a function of:
  - Remaining energy
  - Required transmission energy
- Tries to find traffic distribution (Network flow problem)
- The least cost path is one with the highest residual energy among paths

# Maximum Lifetime Data Gathering



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- Maximizes the Data-gathering schedule
- Maximum Lifetime Data Aggregation
  - Data aggregation plus setting up maximum lifetime of routes
- Maximum Lifetime Data Routing
  - When data aggregation is not possible
- Computational Expensive (scalability)
  - Clustering MLDA



# Minimum Cost Forwarding

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- Aims at finding the minimum cost path in a large network, simple and scalable
- Cost function captures delay, throughput, and energy metrics from node to sink
  - Back-off based algorithm
- Finds optimal cost of all nodes to the sink by using only one message per node
- Does not require addressing or forwarding paths

# Sequential Assignment Routing



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- Table-driven, multi-path protocol
- Creates trees rooted at immediate neighbors of the sink (multiple paths)
  - QoS metrics, energy resource, priority level of each packet
- Failure recoverable (done locally)
- High overhead to maintain tables and states at each sensor



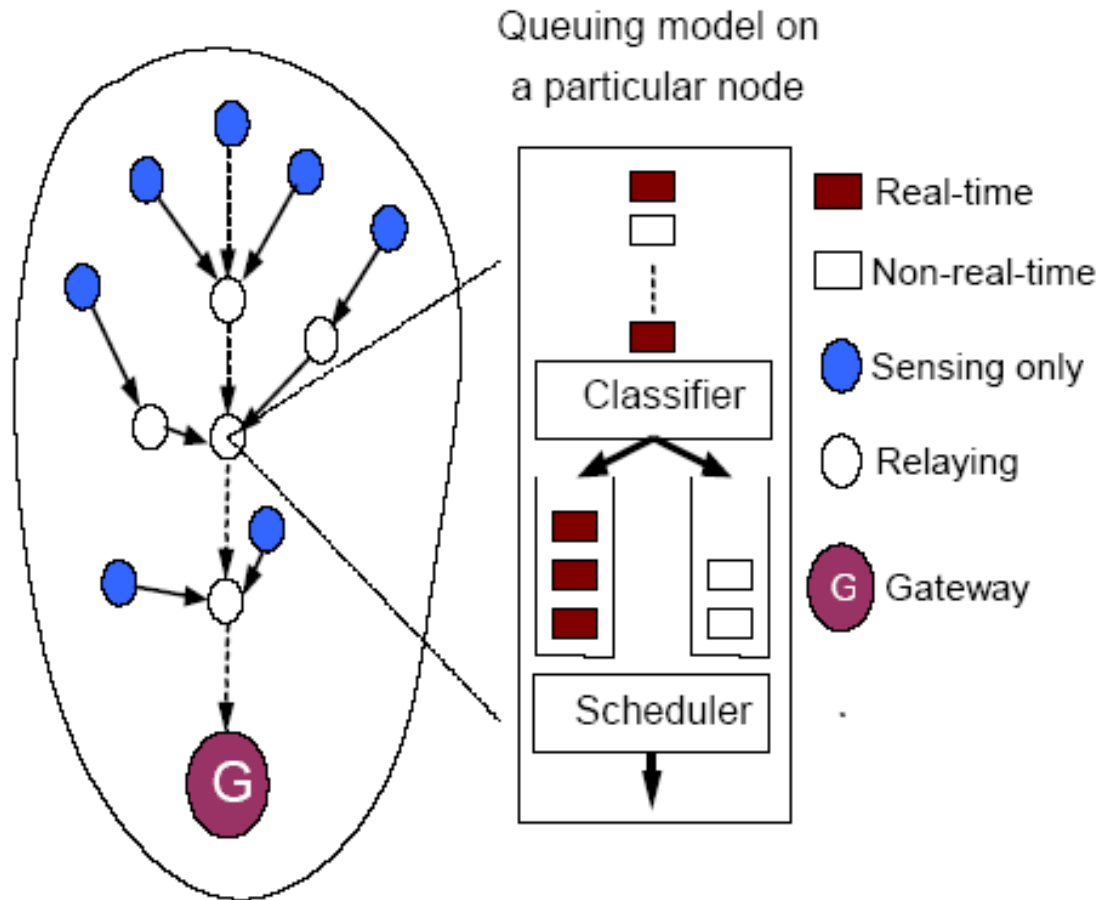
# Energy Aware QoS Routing Protocol



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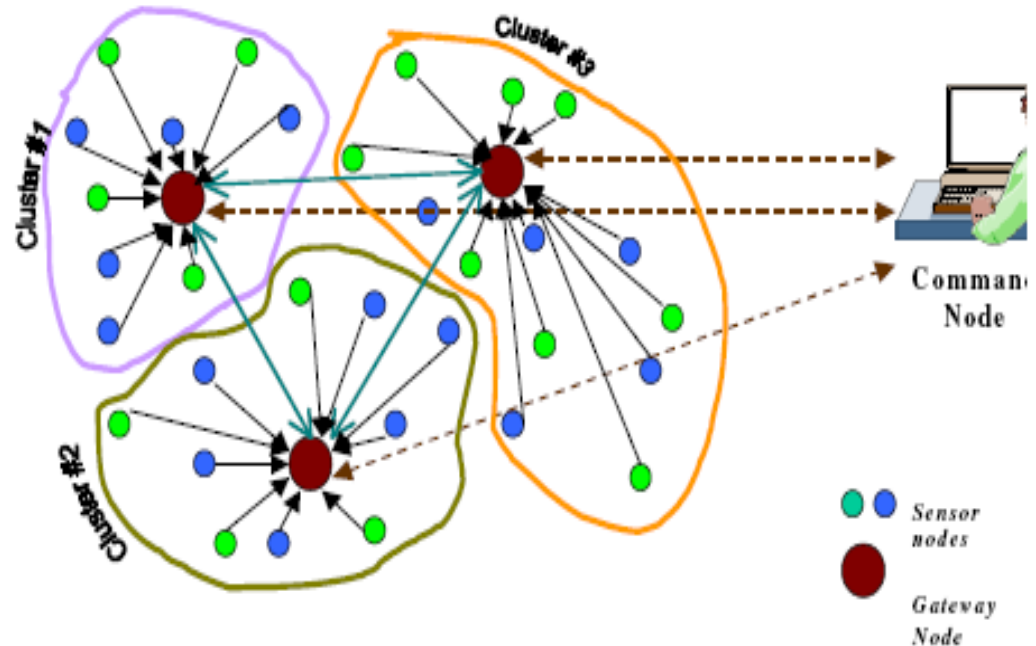
- Finds least cost and energy efficient paths that meet the end-to-end delay during connection
  - Energy reserve, transmission energy, error rate
- Class-based queuing model used to support best-effort and real-time traffic

# Energy Aware QoS Routing Protocol



# Energy Aware QoS Routing Protocol

- Basic settings
  - Base station
  - Gateways can communicate with each other
  - Sensor nodes in a cluster can only be accessed by the gateway managing the cluster
  - Focus on QoS routing in one cluster
  - Real-time & non-real-time traffic exist
    - Support timing constraints for RT
    - Improve throughput of non-RT traffic





# SPEED

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- The protocol requires each node to maintain information about its neighbors and uses geographic forwarding to find the paths.
- SPEED strive to ensure a certain speed for each packet in the network so that each application can estimate the end-to-end delay.
- SPEED can provide congestion avoidance when the network is congested.

# Summary of Routing Protocols in Wireless Sensor Networks

<i>Routing protocol</i>	<i>Data-centric</i>	<i>Hierarchical</i>	<i>Location-based</i>	<i>QoS</i>	<i>Network-flow</i>	<i>Data aggregation</i>
SPIN	✓					✓
Directed Diffusion	✓					✓
Rumor Routing	✓					✓
Shah et al.	✓		✓			
GBR	✓					✓
CADR	✓					
COUGAR	✓					✓
ACQUIRE	✓					
Fe et al.					✓	
LEACH		✓				✓
TEEN&APTEEN	✓	✓				✓
PEGASIS		✓				✓
Younis et al.		✓	✓			
Subramanian et al.		✓				✓
MECN&SMECN			✓			
GAF		✓	✓			
GEAR			✓			
Chang et al.		✓			✓	
Kalpakis et al.			✓		✓	
Akkaya et al.		✓		✓		
SAR				✓		
SPEED			✓	✓		